
Verification of GALS Systems by Combining Synchronous Languages and Process Calculi

Hubert Garavel and Damien Thivolle

INRIA Rhône-Alpes / VASY

<http://www.inrialpes.fr/vasy>



Outline

- **Motivations**
- **TFTP case-study and method**
- **Formal verification**
- **Simulation**



Topic of this talk

- How to perform **model checking** of **GALS**?
- **GALS** (*Globally Asynchronous, Locally Synchronous*)



*synchronous islands in
a sea of asynchrony*



Model checking GALS

- **Why?**
 - **Avionics** companies use synchronous languages:
 - ESTEREL, SCADE, etc.
 - More and more, synchronous components **interact** with an **asynchronous** environment:
 - X-BY-WIRE, Modular Avionics, etc.
- **Our approach:**
 - **Encode** synchronous components as process algebras **functions**
 - Write **wrappers** around the functions for **asynchronous** communications



The CADP toolbox

- A **toolbox** for designing **asynchronous systems**:
 - compilers
 - model checkers
 - equivalence checkers
 - simulation, rapid prototyping, test case generation...
 - performance evaluation
- Developed by the VASY team of INRIA Grenoble
- **43** tools, **18** libraries, **100+** case studies
- **8** supported architectures (32- and 64- bits)
- Important user community (forum...)
- Licensed to several big companies



The TFTP case study



The TFTP case study

- A real example provided to us by Airbus
- Communications between plane and ground
- A three layer protocol stack
 - ARINC protocol
 - TFTP (Trivial File Transfer Protocol)
 - UDP (datagrams over IP)
- Very "light" specification:
 - SAM automaton (7 states, 39 transitions)
 - two TFTP entities connected head-to-tail via UDP

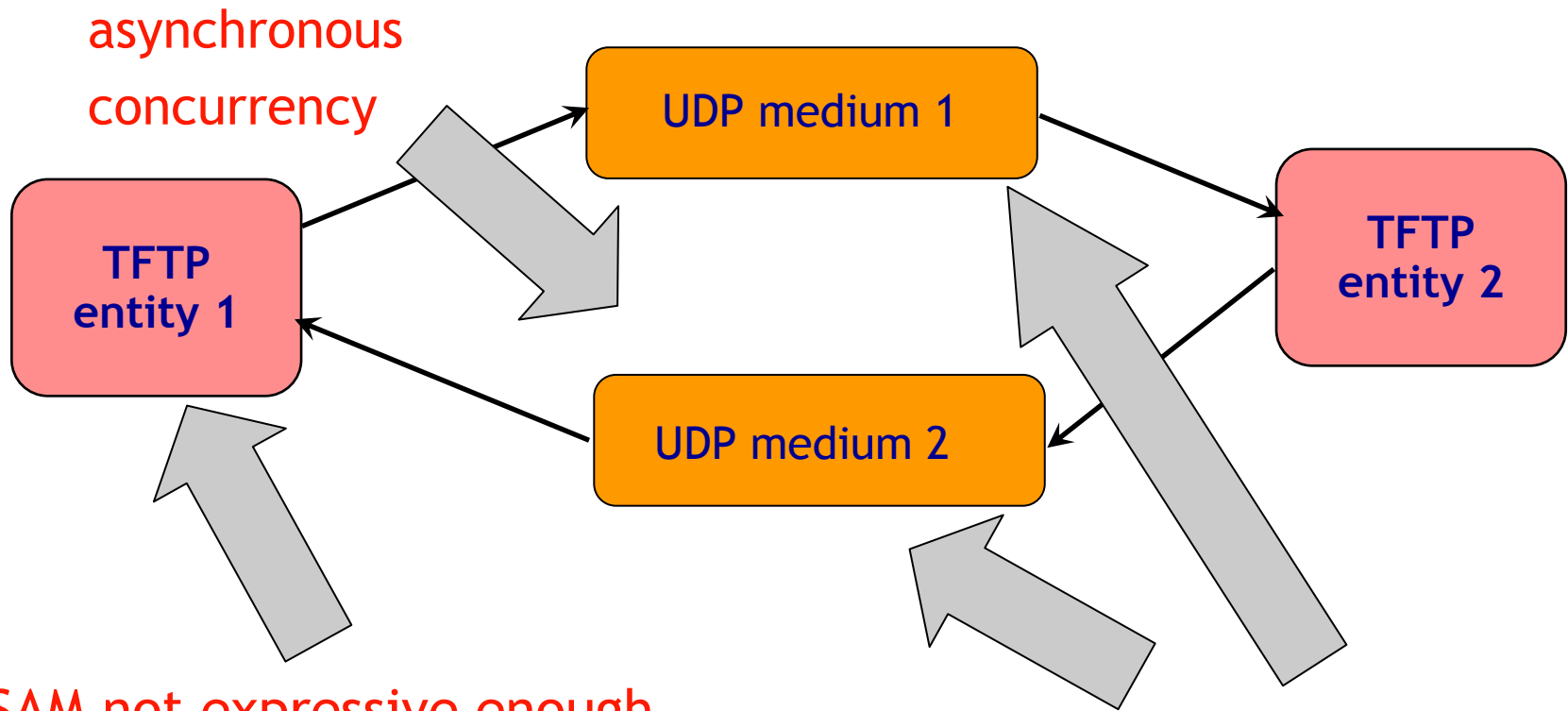


What is SAM?

- A graphical language defined/used by Airbus
- Inspired from F. Maraninchi's Argos language
- A synchronous language:
 - boxes connected by arrows (synchronous parallel, no causality loops)
 - each box is a synchronous automaton
 - boolean inputs/outputs
 - determinism (priority between transitions)
- Reference manual written by VASY:
<http://gforge.enseeiht.fr/docman/view.php/33/2745/SAM.pdf>



SAM limitations as seen on the TFTP



SAM not expressive enough
to describe non-boolean
computations

- message contents
- timeout values

UDP is nondeterministic

- messages can be lost
- message order is not preserved

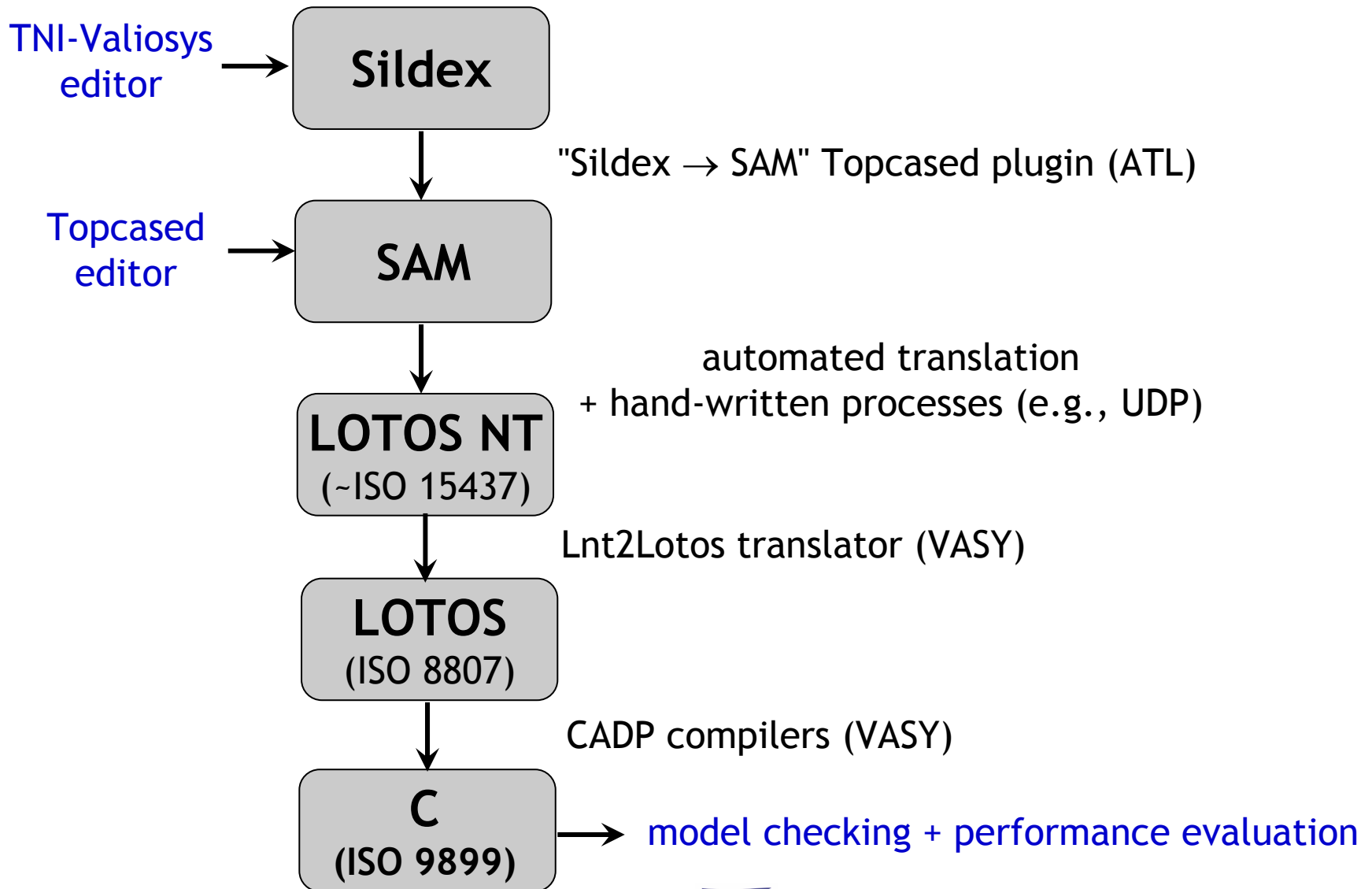


Chosen methodology

- SAM only models a fragment of the problem
- To model and analyze the complete TFTP, we need **an asynchronous language**
- Several attempts made (FIACRE, LOTOS...)
- Best solution chosen: LOTOS NT
 - a subset of the ISO standard E-LOTOS
 - funded and used by Bull



Tool chain



From SAM to LOTOS NT functions

- Synchronous parallelism in SAM \neq
Asynchronous parallelism in FIACRE, LOTOS...
- Each SAM automaton is translated to a sequential Mealy function
 $f(\textit{current_state}, \textit{inputs}) \rightarrow (\textit{next_state}, \textit{outputs})$
- Synchronous composition of SAM automata is implemented by a composition of the corresponding sequential functions
(*in the TFTP, only one SAM automaton*)



From SAM to LOTOS NT functions

function transition (in current:state,
in receive_ACK:bool,
in receive_ERROR:bool,
...
in timeout:bool,
in max_retries_reached:bool,
out next:state,
out send_DATA:bool,
out send_WRQ:bool,
...
out arm_timer:bool,
out stop_timer:bool) is

LOTOS NT:
215 lines
of code

SAM:
7 states
39 transitions

(* Init, out variables := false *)

if current == STATE_6 then

if timeout and not (max_retries_reached) then

send_WRQ := true;
arm_timer := true;
next := current

elsif receive_ERROR then

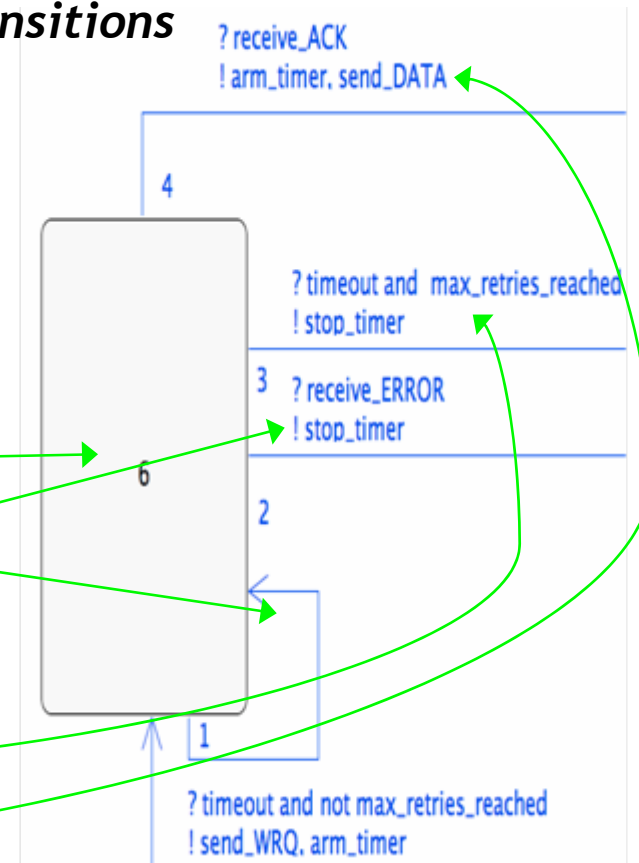
stop_timer := true;
next := STATE_1

elsif timeout and max_retries_reached then

stop_timer := true;
next := STATE_1

elsif receive_ACK then

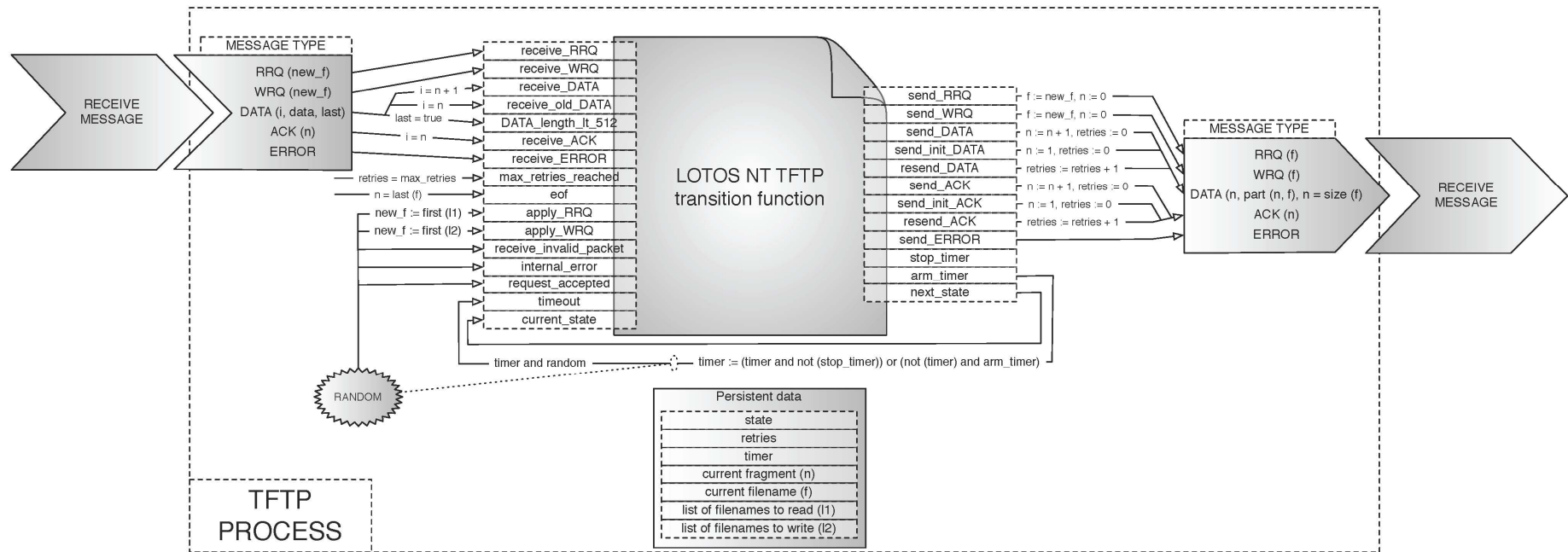
arm_timer := true;
send_DATA := true;
next := STATE_2



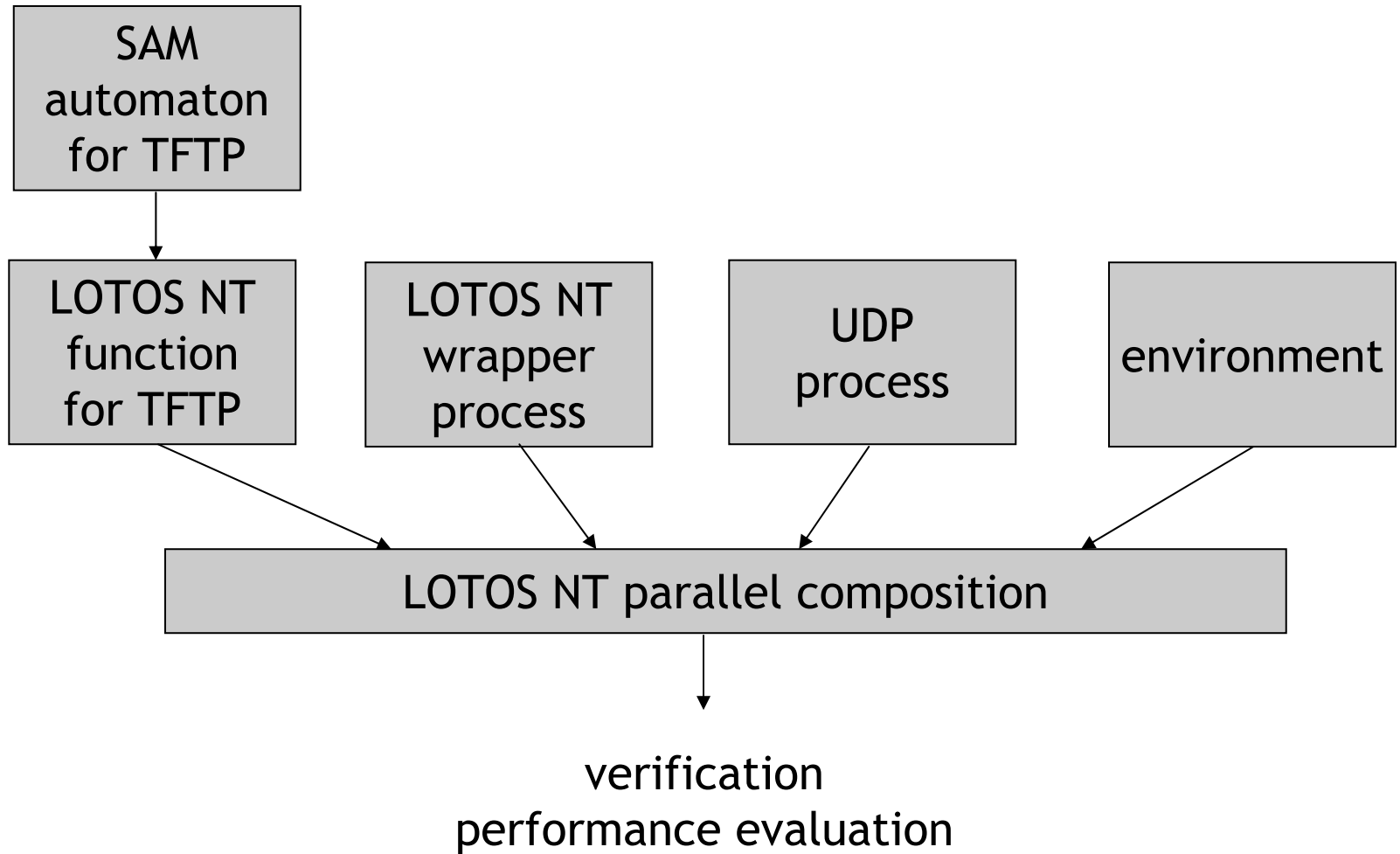
Wrapping functions into LOTOS NT processes

Mealy functions encapsulated into processes:

- converting boolean variables into I/O events
 - Transforms a Mealy function into an LTS
- *adding non-boolean code not described in SAM*



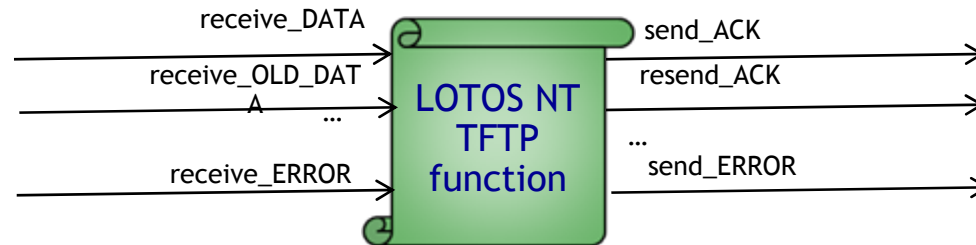
Global view



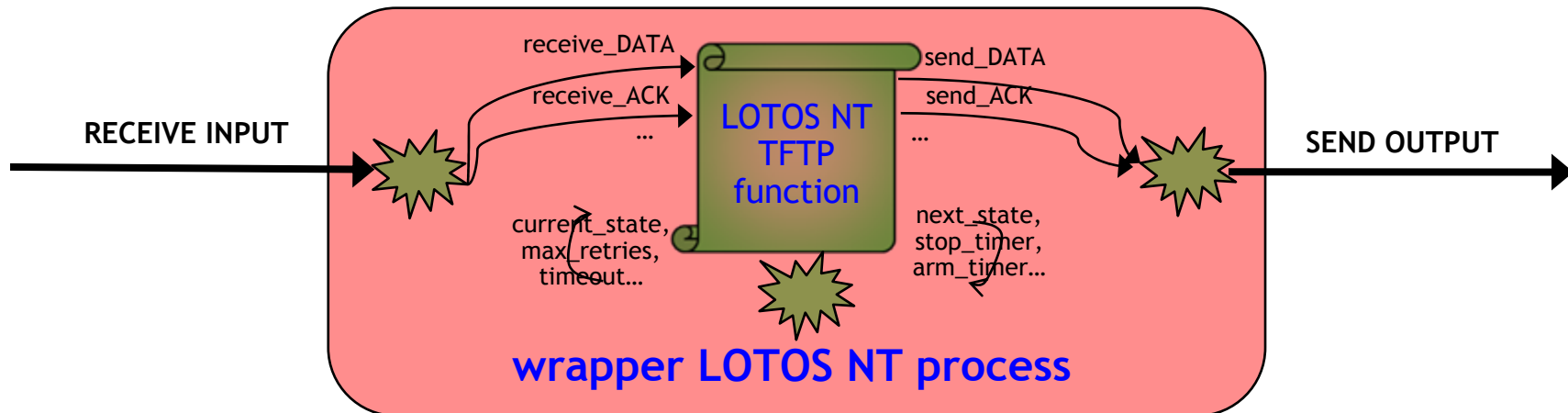
Verification of the "basic" TFTP



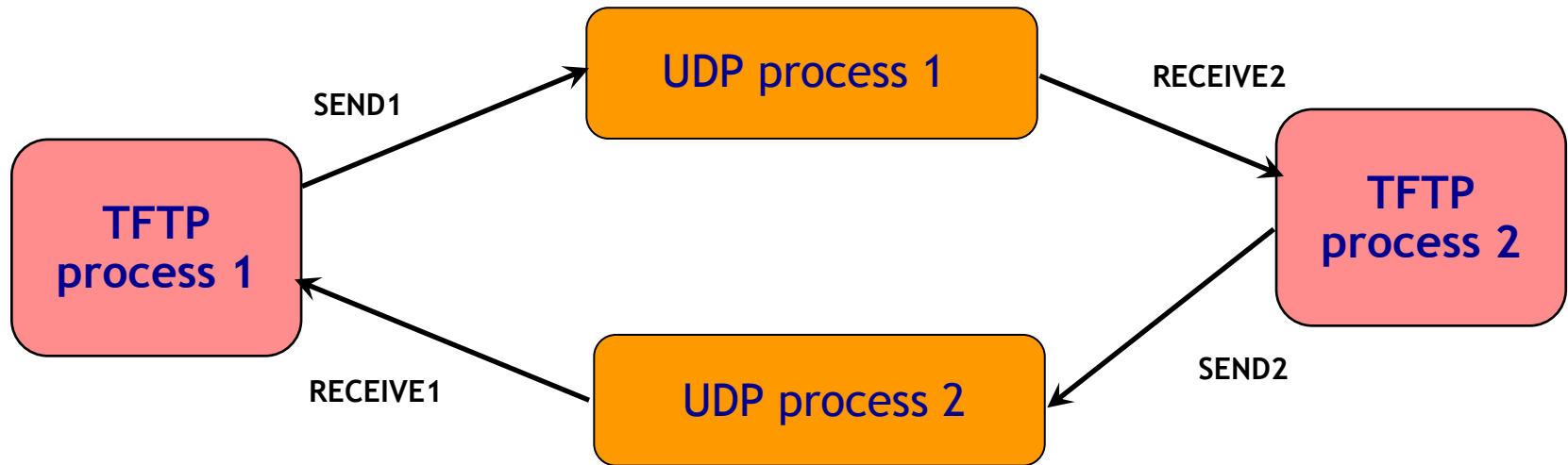
Encapsulation of the SAM automaton



- This function is encapsulated into a wrapper LOTOS NT process
- This wrapper is very simple (193 lines): messages do not carry data



The whole TFTP protocol



- UDP entities are modelled in LOTOS NT too:
 - bounded FIFOs (messages lost)
 - bounded BAGs (messages lost or re-ordered)
- The 4 processes execute asynchronously
- LOTOS NT parallel composition is used for this



State space generation

- Successive steps: LOTOS NT \rightarrow LOTOS \rightarrow LTS
- **Direct** state space generation not efficient
- **Compositional** generation used instead:
 - each sequential process is minimized
 - minimized processes are recombined

Example with two UDP FIFOs of size 2:

- 846,888 states
- 3.7 million transitions
- compositional generation: 15 s



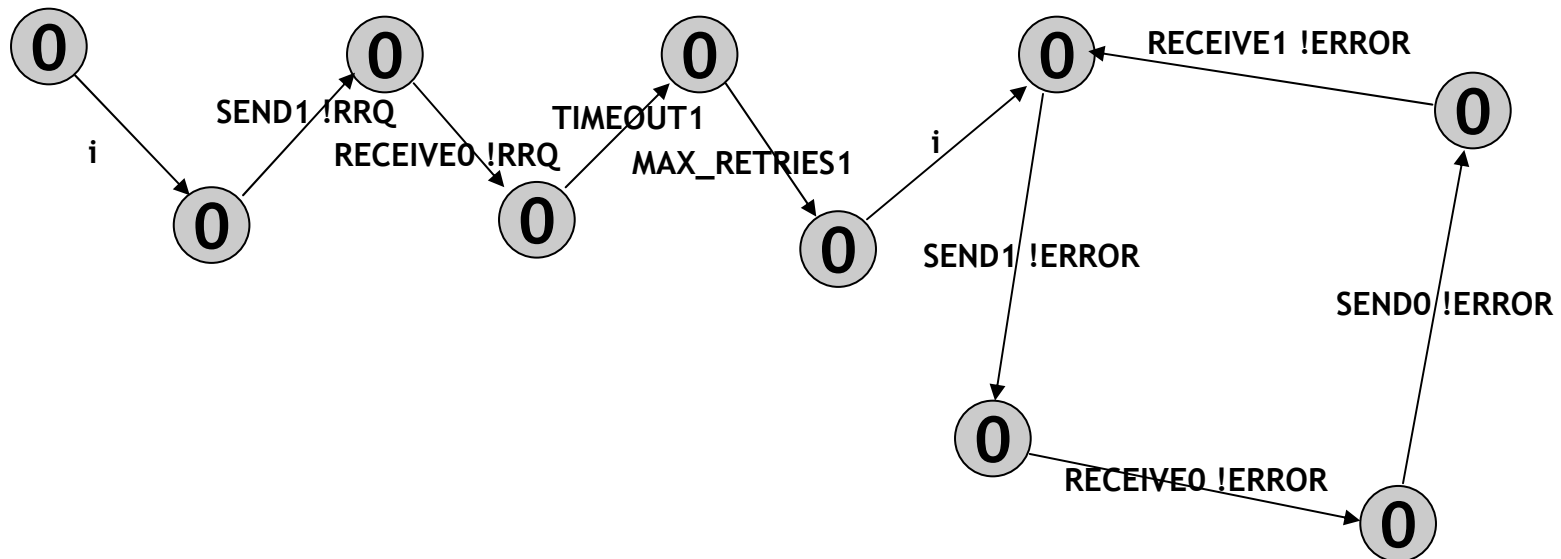
Model checking verification

- 12 properties expressed with Evaluator
- 8 problems detected:
 - timer **does not stop** after a transfer is finished
 - after loss of final ACK, resent DATA is **ignored**
 - new transfer **impossible** right after final ACK
 - invalid packets and invalid acknowledgements are simply **ignored** whereas they should **abort** the transfer
 - ...



Generation of diagnostics

- Property "absence of error loop" not satisfied
- Diagnostic generated by breadth-first search (in 0.53 s):



Verification of the "accurate" TFTP



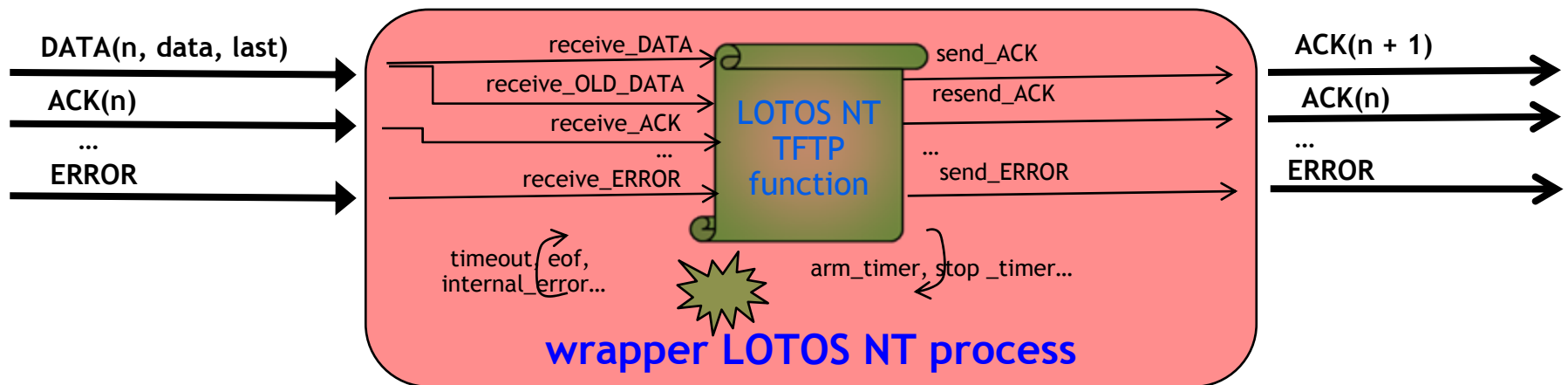
Limitations of the "basic" verification

- The SAM model of the TFTP does not express certain important details:
 - only Boolean variables
 - not represented: counters, number of retries, packet contents, fragment numbers, list of files to be sent or received, etc.
- There are properties that cannot be expressed
 - Example:
ACK (x) cannot be received before ACK (x-1)



"Accurate" TFTP modeling

- The SAM automaton (encoded as a LOTOS NT function) is kept unchanged (215 lines)
- But it is encapsulated in a more elaborate "wrapper" written manually (418 lines) based upon knowledge of the TFTP standard



State space generation

- CADP tools used to generate the state space for various configurations
- Example:
 - TFTP entity 1 has one file to write
 - TFTP entity 2 has one file to read
 - two UDP FIFOs of size 2
 - 44 million states, 221 million transitions
 - compositional generation = 24 mn 22'
- Verified up to size 3 for the UDP FIFOs and BAGs



Results

- 12 formulas of “basic” specification
+ 17 new formulas
- 8 new problems detected:
 - any old DATA received is **not** acknowledged
 - when initiating a write, receipt of an invalid packet is **ignored**
 - when initiating a read, receipt of an invalid packet is **ignored**
 - if both processes send RRQ or WRQ at the **same time**, their requests will **not** be answered
 - INTERNAL_ERROR is **ignored** in several cases
 - etc .



Performance evaluation



Performance issues

- Do problems in TFTP specification affect **runtime performance**?
- If so, how much?
- Example: problem 08 (infinite error loop)
 - in many cases, one may exit this error by reinitializing the TFTP entities after a timeout
 - but timeout reinitialization causes performance degradation
 - can we quantify this degradation?

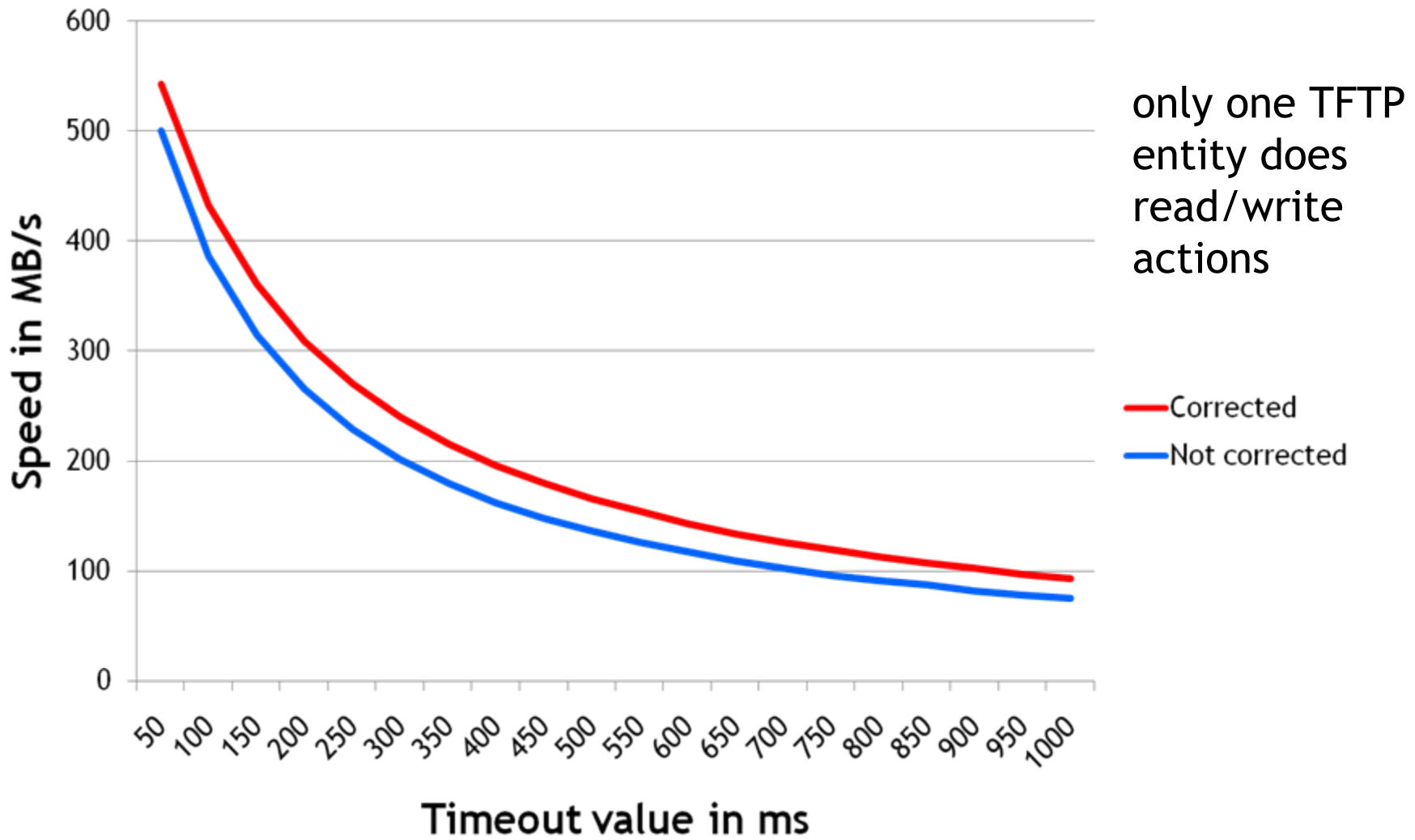


A simulation-based approach

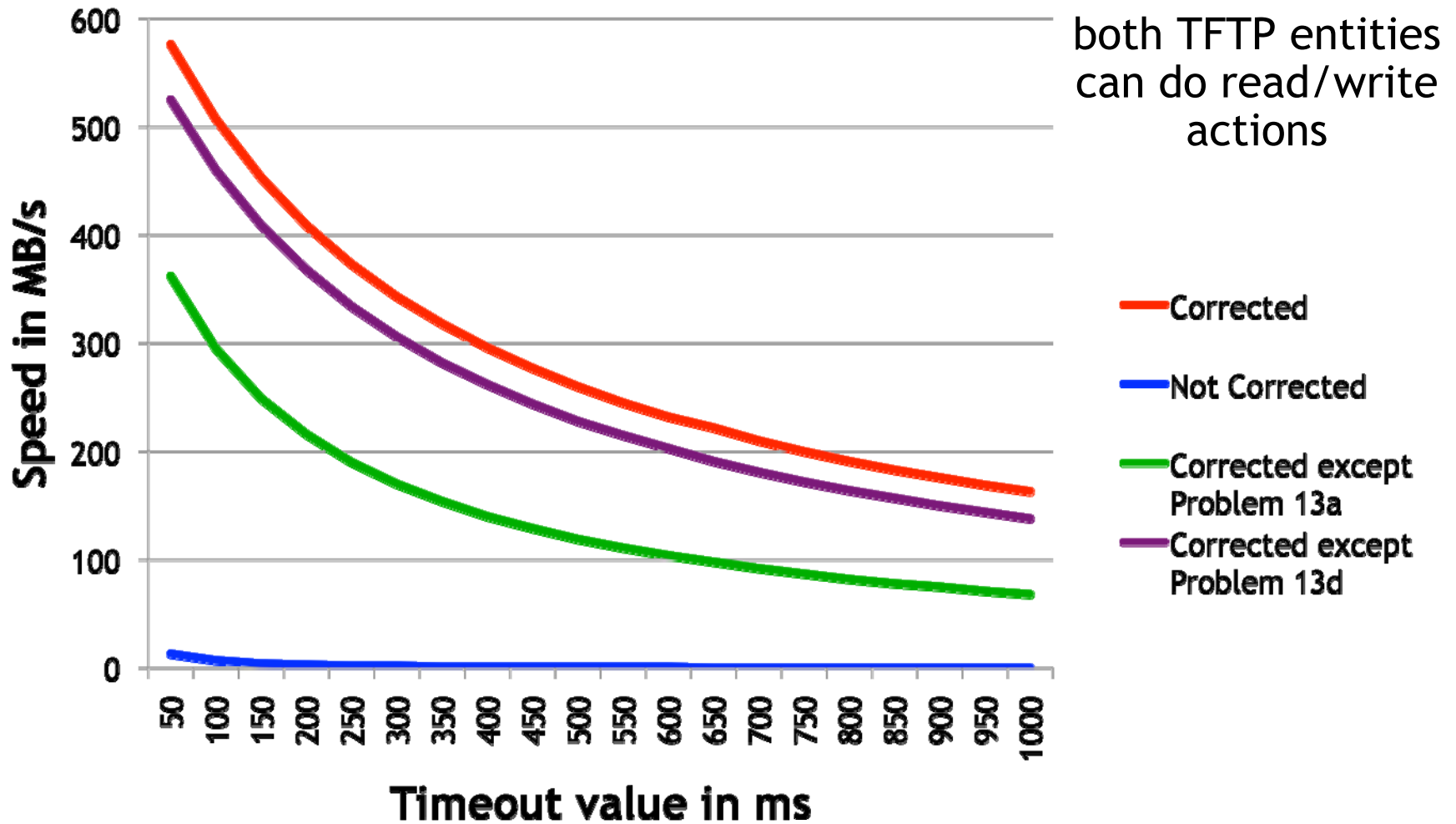
- Instrument the **Executor tool** of CADP
- Generate **random execution traces**
- Measure **TFTP transfer speed** on these traces
- Different scenarios:
 - **scenario 1**: one TFTP entity does read/write
 - **scenario 2**: both TFTP entities do read/write
- Chosen TFTP parameters:
 - **10,000** files written or read in each scenario
 - packet size: **32 kB**
 - medium speed: **1 MB/s**
 - medium latency: **8 ms**
 - medium losses: **1%**



Performance impact of problems (scenario 1)



Performance impact of problems (scenario 2)



Conclusion



Summary of TFTP results

- Verification of the "basic" TFTP:
 - 12 properties checked
 - 8 errors detected
- Verification of the "accurate" TFTP:
 - 29 properties checked
 - 19 errors detected
- Performance evaluation:
 - confirms quantitative issues
 - done by simulation, but other approaches exist in CADP (tools for Markov chains)



Conclusion

- **Model checking of GALS**

- reuse synchronous processes (written in Sildex/SAM) composed asynchronously (in LOTOS NT/LOTOS)
- verification and performance evaluation with CADP
- a sound solution for validating GALS
- mostly automated

- **Positive feedback from Airbus**

- appreciated flow combining Topcased, ATL, and CADP
 - based on formal transformations
 - 5 languages: Sildex → SAM → LOTOS NT → LOTOS → C (+ MCL)
- ongoing collaboration on a new avionics application

