# Testing Resource Isolation for SoC Architectures

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# System-on-Chip & Validation

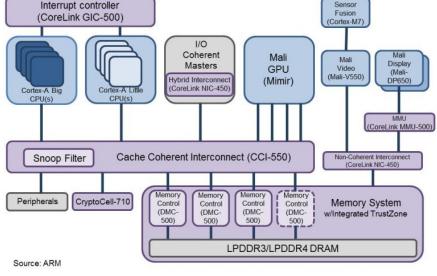
SoC (System-on-Chip) architectures



- Priority: Bug hunting
- Security: **Resource isolation**
- Modern SoCs:

too complex for traditional validation methodologies (directed tests, constrained random test)





# **Model-based System-on-Chip Testing**

- New industrial inclination: Modeling for Testing "Modeling without testing is meaningless"
- **Two** modeling tasks: **behavior** & **test scenario**
- **PSS** (Portable-test and Stimulus Standard)
  - Behavior: *actions* ordered by *flow objects* (buffer, state, stream)
  - Test scenario: verification intent (VI)
     composition of actions with process calculi operators
  - Focus on VI: behavior only to fill gaps in the VI
- Similar to academic conformance testing
  - Generation of test cases for a behavior and test purpose
  - Supported by CADP (LNT language) and TESTOR





#### Outline

Hardware Resource isolation for SoC architectures

Modeling the behavior in LNT and PSS

Modeling the test scenarios and generating tests



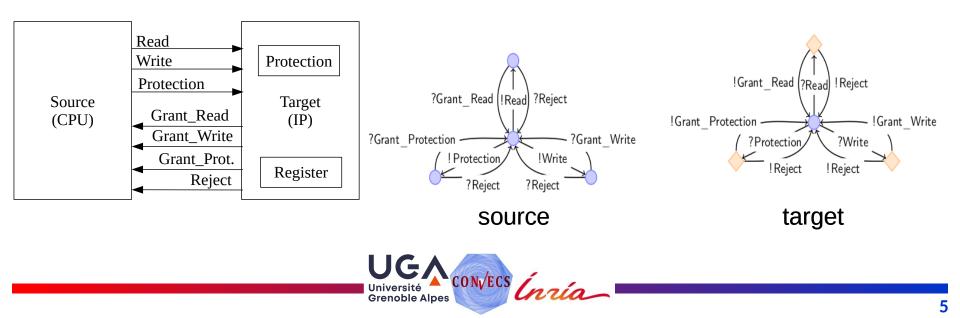


### **Hardware Resource Isolation**

- Mechanism to ensure a program or IP cannot access data or functionalities not intended for it
  - ARM PSA (Platform Security Architecture)



- Security: Secure/Non-secure (TrustZone)
- Privilege: Privileged/Non-privileged (elevation levels EL<sub>0</sub>-EL<sub>3</sub>)



### LNT code for TARGET

```
process TARGET [Read, Grant Read, Reject Read, Write, Grant Write,
1
                    Reject Write, Protection, Grant Protection,
2
3
                    Reject Protection: Bus] (id: ip) is
4
      require not (source (id));
5
      var d,e: data, s,t,u: security, p,q,r: privilege, o, other: ip in
         d := data1; --- default value
6
7
          s := non secure; p := non privileged; --- lowest protection level
8
         loop
9
             select
                Read (?o, id, ?t, ?q) where source (o);
10
                if valid access (s, t, p, q) then
11
                   Grant_Read (o, id, d)
12
13
                else
                   Reject Read (o, id)
14
                end if
15
. . .
            -- communication between other IPs on the shared interconnect
30
31
             [] Read (?other, ?o, ?any security, ?any privilege)
32
                   where (o != id) and source (other)
. . .
                                1 process LNT per IP
50
             end select
51
         end loop
52
      end var
                                Rendezvous on the same gates
53
   end process
                                                                      (actions)
                                       CONVECS
```

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# **LNT Behavior Modeling Results**

Several equivalent models (hiding source IDs)

- 8 sources (stable configuration) and 1 target
   182 states, 558 transitions, and 99 labels
- 1 source (changing configuration) and 1 target
   52 states, 268 transitions, and 39 labels
- Model checking of temporal logic properties (e.g., each request is followed by a response, illegal requests are rejected, ...)
- Large state spaces for more than 1 target



## **PSS Behavior Modeling**

#### Inspired by the 1 source/1 target LNT model

- 21 actions
- 2 state FOs (source and target)
- 9 stream FOs (to emulate rendezvous)
- constraints
   to indicate
   unchanged
   state fields

```
action t request read {
 input target state in state;
         request read stream in stream;
 input
 output target state out state;
  constraint in state.initial == false;
  // Idle \rightarrow Read
  constraint in state.sstate == idle;
  constraint out state.sstate == read;
  // save stream data
  constraint out state.tx sec
                                     = in stream.sec;
  constraint out state.tx priv
                                     = in stream.priv;
  // Maintain fields
  constraint out state.data
                                  = in state.data;
  constraint out state.sec
                                  = in state.sec;
  constraint out state.priv
                                  = in state.priv;
  constraint out state.tx data
                                  = in state.tx data;
  constraint out state.next sec
                                  = in state.next sec;
  constraint out state.next priv == in state.next priv;
```

Tedious, error-prone, > 500 lines, huge state space 1.7 billion states, 14 billion transitions, 7000 labels

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### **Monolithic PSS Behavior Modeling**

Monolithic, a complex	<pre>ction t_grant_read {     input system_state in_state;     output system_state out_state;</pre>
state	<pre>constraint in_state.initial == false; // Move from Read to Idle</pre>
No streams	<pre>constraint in_state.sstate == read; constraint out_state.sstate == idle;</pre>
11 actions	<pre>// Check protection constraint (in_state.source_sec == secure)    (in_state.target_sec == non_secure);</pre>
Less modular	<pre>(in_state.target_sec == non_secure); constraint (in_state.source_priv == privileged)    (in_state.target_priv == non_privileged);</pre>
More	<pre>// Maintain source fields constraint out_state.source_sec == in_state.source_sec;</pre>
constraints	<pre>constraint out_state.source_priv == in_state.source_priv; constraint out_state.source_data == in_state.source_data;</pre>
Bisimilar to	<pre>// Maintain target fields constraint out_state.target_sec == in_state.target_sec;</pre>
LNT model	<pre>constraint out_state.target_priv == in_state.target_priv; constraint out_state.target_data == in_state.target_data;</pre>
(after renaming	
and hiding) }	UCA Université Grenoble Alpes

### **Test Generation from Test Scenarios**

#### Test scenario:

partial ordering of some actions from the behavior

- Two test scenarios illustrating both methodologies (two more test scenarios in the paper)
  - Differences of the methodologies:

	PSS methodology	Conformance testing
Test scenario	Verification intent	Test purpose
<b>Test generation</b>	Backward inference	Forward exploration



### **Test 2: Interleaving of all Responses**

1	process PURPOSE_2 [ LNT	acti
2	Reject Read,	t
3	Reject Write ,	
4 5	Reject Protection ,	t
5	Grant_Read,	t
6	Grant_Write ,	t
7	Grant_Protection ,	t t t t
8	TESTOR_ACCEPT: none] is	а
9	par	
10	Grant_Read	
11	Grant_Write	
12	Grant_Protection	
13	Reject Read	
14	Reject Write	
15	Reject Protection	
16	end par;	
17	loop TESTOR_ACCEPT end loop	}
18	end process	}
- <u>-</u>		•

```
PSS
ion intent 2 {
grant read
                   Grant Read;
                   Grant_Write;
grant write
_grant_protection Grant_Protection;
reject read
                   Reject Read;
reject write Reject Write;
_reject_protection Reject_Protection;
activity {
 schedule{
     Grant Read;
     Grant Write:
     Grant Protection;
     Reject Read;
     Reject Write;
     Reject Protection;
 }
```

PSS: only shortest tests without repetitions

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LNT: all tests with coverage guarantees

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#### **Test 4: Access data with different protection**

```
process PURPOSE 4 [Read, Grant Read, Write, Grant Protection: Bus,
 1
                         TESTOR ACCEPT, TESTOR REFUSE: none] is
2
       var s,t: security, p,q: privilege, d: data in
3
          Grant Protection (?any ip, ip0, ?s, ?p)
4
          Write (?any ip, ip0, s, p, ?d); — same s and p as in the previous line
5
          select
6
7
             — refuse any further rendezvous on gate Grant_Protection
8
              Grant_Protection (?any ip, ip0, ?s, ?p); loop TESTOR REFUSE end loop
9
          [] — accept all other rendezvous
              null
10
          end select:
11
          Read (?any ip, ip0, ?t, ?q) where (s != t) or (p != q);
12
13
          Grant Read (?any ip, ip0, d); — access data with different security and privilege levels
          loop TESTOR ACCEPT end loop
14
15
       end var
16
   end process
```

#### Cumbersome and error-prone to express in PSS



# Conclusion

This talk: Compare modeling & testing approaches of PSS and LNT

Formal modeling in the hardware design domain

Over the second the second the second the second test generation is considered test generation.

Building complete system models is not envisaged

PSS enables modeling in view of test generation but does *not* enable conformance testing

Perspective: Combine both worlds

- formal model-based conformance testing as front-end
- PSS test execution as back-end



#### **Thank You**

For Further Information



#### https://accellera.org/down loads/standards/portablestimulus



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https://cadp.inria.fr